

$$X = \{0, 1\} \quad Y = \{0, 1\} \quad \beta: X^+ \rightarrow Y$$

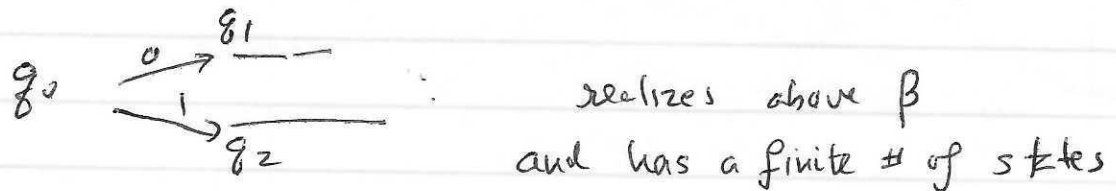
$$\beta(x) = \begin{cases} 1 & \text{if } \# \text{ of } 1\text{'s in } x \\ & \text{greater or equal to } \# \text{ of } 0\text{'s.} \\ 0 & \text{otherwise} \end{cases}$$

$$\# 1\text{'s} \geq \# 0\text{'s}$$

Example

10110	output is 1
00111	output is 1
001	output is 0
0000011110	output is 0.

Can we build such a finite-state Mealy Machine



First: What is the intuition?

Second: How do you prove this

$\beta(z)$ is the behavior (starting)

$$\beta \circ h_x(z) = \beta(xz)$$

Goal choose a set of behavior functions that are different $\hat{=}$ an countably infinite number.

$|\{ \beta \circ h_x \mid x \in X^* \}|$ is not finite.

z	$\beta(z)$	$\beta \circ \beta_0(z)$	$\beta \circ \beta_{00}(z)$	$\beta \circ \beta_{000}(z)$	$\beta \circ \beta_{0000}(z)$...
1	1	1	0	0	0	
11	1	1	1	0	0	
111	1	1	1	1	0	
1111	1	1	1	1	1	
⋮						

We illustrate that \exists an infinite # of behavior functions for this $\beta: X^+ \rightarrow Y$ function. Hence it is not finite state realizable

State Minimization Algorithm

Example Moore Machine

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First partition on output function λ .
 class 1 $\{0, 3\}$ class 2 $\{1, 2, 4, 5\}$

		Q \ X	
		0	1
class 1	0	0: c1	1: c2
	3	3: c1	4: c2
class 2	1	1: c2	2: c2
	2	2: c2	3: c1
	4	4: c2	5: c2
	5	5: c2	0: c1

Note no change to class 1, cannot be further split at this time
 class 2 must be split
 $\{1, 4\} \approx \{2, 5\}$

		Q \ X	
		0	1
class 1	0	0: c1	1: c2
	3	3: c1	4: c2
class 2	1	1: c2	2: c3
	4	4: c2	5: c3
class 3	2	2: c3	3: c1
	5	5: c3	0: c1

No change to any classes
 Algorithm terminates

Equivalence sets of states are:

$\{0, 3\}$, $\{1, 4\}$, $\{2, 5\}$

Reduced machine has 3 states.