Introduction to the Domain Name Service (DNS)

- 1. Host name to IP number mapping was originally done by downloading a static file
- 2. The UNIX version of this file is /etc/hosts (the file we ftp'd in was called hosts.txt



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3. The central file was maintained by the Stanford Research Institute Network Information Center (SRI-NIC)



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4. As the Internet grew this scheme became unworkable
The size of the file became too large
The load on SRI-NIC site became too heavy
The file was always inconsistent with reality
Hostname collisions became frequent (anyone could name their machine "whitehouse.gov" if they wanted to)



Overview of DNS

In 1984 Paulk Mockapetris of USC designed the architecture of DNS. It's based on the idea of "resource records".



Overview of DNS

The InterNIC was set up to manage DNS; however, this responsibility was given over to ICANN (Internet Corporation for Assigned Names and Numbers), and now the "InterNIC" is just a website. One of the most important activities of ICANN is to accredit registrars.



Overview of DNS

Today, there are hundreds of registrars, and prices for domain names have dropped under \$10.



Features

Is Local control: each segment is updated locally

Global access: each segment is available (almost) immediately to the rest of the world upon update



Features

Robustness: achieved through replication

Real Adequate performance: is achieved through caching



Software

Servers: called name servers, contain information about some segment of the network and make it available to clients ("BIND" = "Berkeley Internet Name Daemon", includes "named", libraries, "nslookup", "dig", "host")



Software

Client: resolvers, a set of library routines that resolve names by accessing a server (originally a separate library, like libresolv.a, now usually part of libc.a)



Software

Domain name server software is also available for non-UNIX platforms, such as Windows 2000/2003 and Macintosh OS X.



Similar to the structure of a hierarchical file system

The root's name is the null label "" but is written as a single dot "."



Reach node represents a 'domain'

Regional Every domain is named



The full domain name is the sequence of labels from the domain to the root, separated by periods



Image: Unlike a file system pathname the name is read from leaf to root (right to left rather than left to right)

xi.cs.fsu.edu



Domain management

- Each domain may be managed by a different organization
- The organization may divide itself into subdomains



Domain management

Image: Then delegate responsibility for maintaining them

ICANN (currently) manages the "provisioning" of toplevel domains



Domain management

What is ICANN?

The Internet Corporation for Assigned Names and Numbers (ICANN) is responsible for managing and coordinating the Domain Name System (DNS) to ensure that every address is unique and that all users of the Internet can find all valid addresses. It does this by overseeing the distribution of unique IP addresses and domain names. It also ensures that each domain name maps to the correct IP address.

ICANN is also responsible for accrediting the domain name registrars. "Accredit" means to identify and set minimum standards for the performance of registration functions, to recognize persons or



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entities meeting those standards, and to enter into an accreditation agreement that sets forth the rules and procedures applicable to the provision of Registrar Services.



Host names

Reach host on a network has a domain

The domain points to information about the host



Host names

This may include:

- An IP address (A records)
- Mail routing information (MX records)



Host names

Solution → Aliases which point to the real ("canonical") host name (CNAME records)



Image There may be any number of branches at a node

Some implementations limit the tree's depth



Image Teach name may contain up to 63 characters

INSIGN The suggested length is 12 or less characters



- A domain name that is written relative to the root is called a 'fully-qualified domain name' - FQDN
- Names without trailing dots ("leading dots") are sometimes interpreted as relative to some domain other than root



Sibling nodes must have unique names

The name of a domain is the domain name of the node at the top of the domain (example purdue.edu)



Reading Again, similar to a file system

Real A node is in multiple domains



- So, a domain is just a subtree of the domain name space ("subdomain")
- Must not use "_", although other similiar naming schemes (prominently, such as NIS) have allowed this.



INST Where are the hosts?

Real A domain name is just an index into the DNS database



The 'hosts' are domain names that point to individual machine information

The hosts are related 'logically' usually by geography or organization



INFIGURE They are NOT necessarily related by network or IP address or hardware type

You could have 10 different hosts on 10 different networks in ten different countries all in the same domain (hp.com)



Nodes at the leaves of the tree usually represent individual hosts

Interior nodes may point to both host information and to subdomain information For example, "hp.com" is both the name of a domain and the name of a machine that routes mail



r Terms

- >>> top-level domain (TLD): a child of root (edu)
- ➡→ first-level domain: a child of root (edu)
- Second-level domain: a child of 1st level domain (fsu.edu)





Reveal Naming rules - the original 7 top-level domains were:

- ➤→ commerical organizations
- edu educational organizations
- >>> gov governmental bodies
- mil military organizations
- met networking organizations
- >> org non-commercial organizations
- >> int international organizations



International names (ISO 3166-1 names)

2-letter designations are reserved for each country (e.g.: DE - Germany, DK - Denmark, CH - Switzerland)
 Each country may organize its domain space however it wishes


The domain name space

- >>> For example, Australia uses edu.au and com.au
- And Britain uses co.uk corporations, ac.uk academic community
- >> And the U.S. uses states: fl.us, then cities: tlh.fl.us



Ref Zones

 A program that stores information about the domain name space is called a Domain Name Server
 A name server generally has complete information about some part of the domain name space



- The subspace is called a 'zone'
- The server is said to have 'authority' for one or more zones
- ≫ What is the difference between a zone and a domain?
 [A domain may be composed of one or more zones, but not vice versa]



- Image: Types of name servers



Secondary master

- Gets the data for its zone from another server
- It periodically updates its local data by copying the primary master's files
- This is called a 'zone transfer'



Solution → Generally keep more than one name server for any given zone (1) Redundancy: fault tolerance (2) Load: localize it as much as possible



Name service clients

Image These are the clients that access name servers

In BIND these are a set of library routines In BIND these are a set of library routines



Name service clients

These are compiled (or linked via shared library) into ssh, sftp, scp, telnet, etc. so that these programs will use DNS to resolve names ("gethostbyname()" and others)



Duties of a simple resolver

- Sometimes called a 'stub resolver'
- Real Querying a name server
- Interpreting the response
- Resend a response
- Returning a reply to the program that it is servicing



How does the name server resolve names

If the name is in the name server's zone then it can give the resolver an immediate 'authoritative' response



How does the name server resolve names

If not, then the name server must search the domain name space for an answer



- Technical operation information about the root servers can be found at http://www.root-servers.org/
- Is The root name servers are authoritative for the toplevel domains (edu, org, us, dk, etc.)



Operator	Location	IP	
Verisign	VA	198.41.0.4	
ISI	CA	192.228.79.201	
		2001:478:65::53	



Operator	Location	IP
Cogent	VA; CA; NY; IL	192.33.4.12
UMD	MD	128.8.10.90
NASA	CA	192.203.230.10



Operator	Location	IP
ISC	Canada; CA; NY;	192.5.5.241
	Spain; Hong Kong	2001:500::1035
	Italy; NZ	
	Brazil; China	
	SK; Russia; Taipei	
	Dubai; France;	
	Singapore	
	Portugal; ZA	



Operator	Location	IP
ISC (cont'd)	Tel Aviv; Jakarta	
	Germany; Japan	
	CZ; Netherlands	
	Kenya; India	
	Britain; Chile	
	Pakistan	



Operator	Location	IP
DOD NIC	VA	192.112.36.4
Army	MD	128.63.2.53
		2001:500:1::803f:235



They can point you to the name servers for each of the top-level domains

They, in turn can point you to their subdomains, etc. until the name is resolved



This scheme puts a lot of importance on the root-level servers



Recursion

The first name server can make multiple requests

- Successive requests refer the first server to another machine
- Real A local server generally responds to a 'recursive query'



Recursion

This is because the local 'dumb' resolver is not smart enough to follow any referrals

Real A recursive query places most of the work on a single name server



Recursion

Image When a recursive query is made the name server is obliged to go find the answer or return an error message



- What if you have an IP number and want to find the host name?
 - Solution → This is useful to make output more readable
 → Used for some security checks



This was easy with the old /etc/hosts tables

Image: The DNS data is indexed by name



- Could do an exhaustive search
- ⇒ The clever solution



Create a part of the domain name space that uses addresses as names



For example type:

nslookup 128.186.120.2 Server: 128.186.120.179

Address: 128.186.120.179#53

2.120.186.128.in-addr.arpa name = diablo.cs.fsu.edu.



Now, as you can see, newer **nslookup** versions will do this automatically. (Of course, you get many versions today of **nslookup** that announce that nslookup is "deprecated".)



Each time a local name server processes a recursive query it learns a lot of information

Image This is cached which speeds up successive queries



Rear Example:

- >>> Say our server has already looked up the address of eecs.berkeley.edu
- This means it has cached the name servers for both eecs.berkeley.edu and berkeley.edu
- >>> If we now make a query for baobab.cs.berkeley.edu
 the local server can skip the root-level query and go right
 to berkeley.edu



Image: Time to live (TTL)

TTL is the amount of time that information is cached before it is discarded



Image The trade-off is between consistency and performance

Remember, caching is also performed by other actors: in particular, **nscd** can cache names, and it uses its *own* TTLs, not those of the actual RRs. Applications such as browsers often also do a good bit of caching of ip values.



Configuring DNS: Client side

Setting up clients



Overview of DNS

The client will try the nameservers in order: "nu", then "mailer", then "trantor"

- you can comment out nu and/or mailer then use nslookup and see results
- INF or put a bogus address in the first entry to see if the resolver tries number 2

Image: The changes take effect immediately



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nslookup chi Server: TRANTOR.UMD.EDU Address: 128.8.10.14

- Name: chi.cs.fsu.edu
- Address: 128.186.121.20



The named.conf file

```
11
// named.conf for Red Hat caching-nameserver
//
options {
        directory "/var/named";
        dump-file "/var/named/data/cache_dump.db";
        statistics-file "/var/named/data/named_stats.txt";
        /*
         * If there is a firewall between you and nameservers you want
         * to talk to, you might need to uncomment the query-source
         * directive below. Previous versions of BIND always asked
         * questions using port 53, but BIND 8.1 uses an unprivileged
         * port by default.
```


```
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```

```
*/
         // query-source address * port 53;
};
11
// a caching only nameserver config
11
controls {
        inet 127.0.0.1 allow { localhost; } keys { rndckey; };
};
zone "." IN {
        type hint;
        file "named.ca";
};
zone "localdomain" IN {
        type master;
        file "localdomain.zone";
        allow-update { none; };
```



};

```
zone "localhost" IN {
    type master;
    file "localhost.zone";
    allow-update { none; };
};
```

```
zone "0.0.127.in-addr.arpa" IN {
    type master;
    file "named.local";
    allow-update { none; };
};
```



```
zone "255.in-addr.arpa" IN {
    type master;
    file "named.broadcast";
    allow-update { none; };
};
```

```
zone "0.in-addr.arpa" IN {
    type master;
    file "named.zero";
    allow-update { none; };
};
```

include "/etc/rndc.key";



Setting up a caching-only server: used to be more popular, now **nscd** is more popular

Still it is very easy to do these days: yum y caching-nameserver, then just turn on default installation /etc/init.d/named start and change /etc/resolv.conf

[root@sophie root]# nslookup

> www.yahoo.com

Server: 127.0.0.1



Address: 127.0.0.1#53

```
Non-authoritative answer:
www.yahoo.com canonical name = www.yahoo.akadns.net.
Name:
     www.yahoo.akadns.net
Address: 68.142.226.43
     www.yahoo.akadns.net
Name:
Address: 68.142.226.45
Name: www.yahoo.akadns.net
Address: 68.142.226.50
     www.yahoo.akadns.net
Name:
Address: 68.142.226.35
Name:
     www.yahoo.akadns.net
```



Address: 68.142.226.38 Name: www.yahoo.akadns.net Address: 68.142.226.39 Name: www.yahoo.akadns.net Address: 68.142.226.41 Name: www.yahoo.akadns.net Address: 68.142.226.42



Logging and named

errors: like most daemons, **named** errors (and other information) are routed through syslog, which you control wtih /etc/syslog.conf:

Log all kernel messages to the console.
Logging much else clutters up the screen.
#kern.*

/dev/console

Log anything (except mail) of level info or higher. # Don't log private authentication messages! *.info;mail.none;news.none;authpriv.none;cron.none

/var/log/messages



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The authpriv file has restricted access.
authpriv.*

Log all the mail messages in one place.
mail.*

#	Log	cron	stuff			
cron.*						

#	Everybody	gets	emergency	messages
*.	emerg			

Save news errors of level crit and higher in a special file. uucp,news.crit /var/log/spooler

Save boot messages also to boot.log
local7.*



*





/var/log/secure

/var/log/maillog

/var/log/cron

INN
#
news.=crit
news.=err
news.notice

/var/log/news/news.crit
/var/log/news/news.err
/var/log/news/news.notice

And here is what you see in /var/log/messages



Feb 14 10:18:20 sophie named[7597]: zone localdomain/IN: loaded serial 42
Feb 14 10:18:20 sophie named[7597]: zone localhost/IN: loaded serial 42
Feb 14 10:18:20 sophie named[7597]: running



DNS Security

INFO fields

- Image: "basic": limit zone transfers
- $\ensuremath{\mathbb{R}}\xspace^{\ensuremath{\mathbb{R}}\xspace}$ use keys and $\ensuremath{\textbf{rndc}}\xspace$

