

CIS 5371, FALL 2024

SYMMETRIC ENCRYPTION

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The slides are loosely based on those of
Prof. Mihir Bellare, UC San Diego.

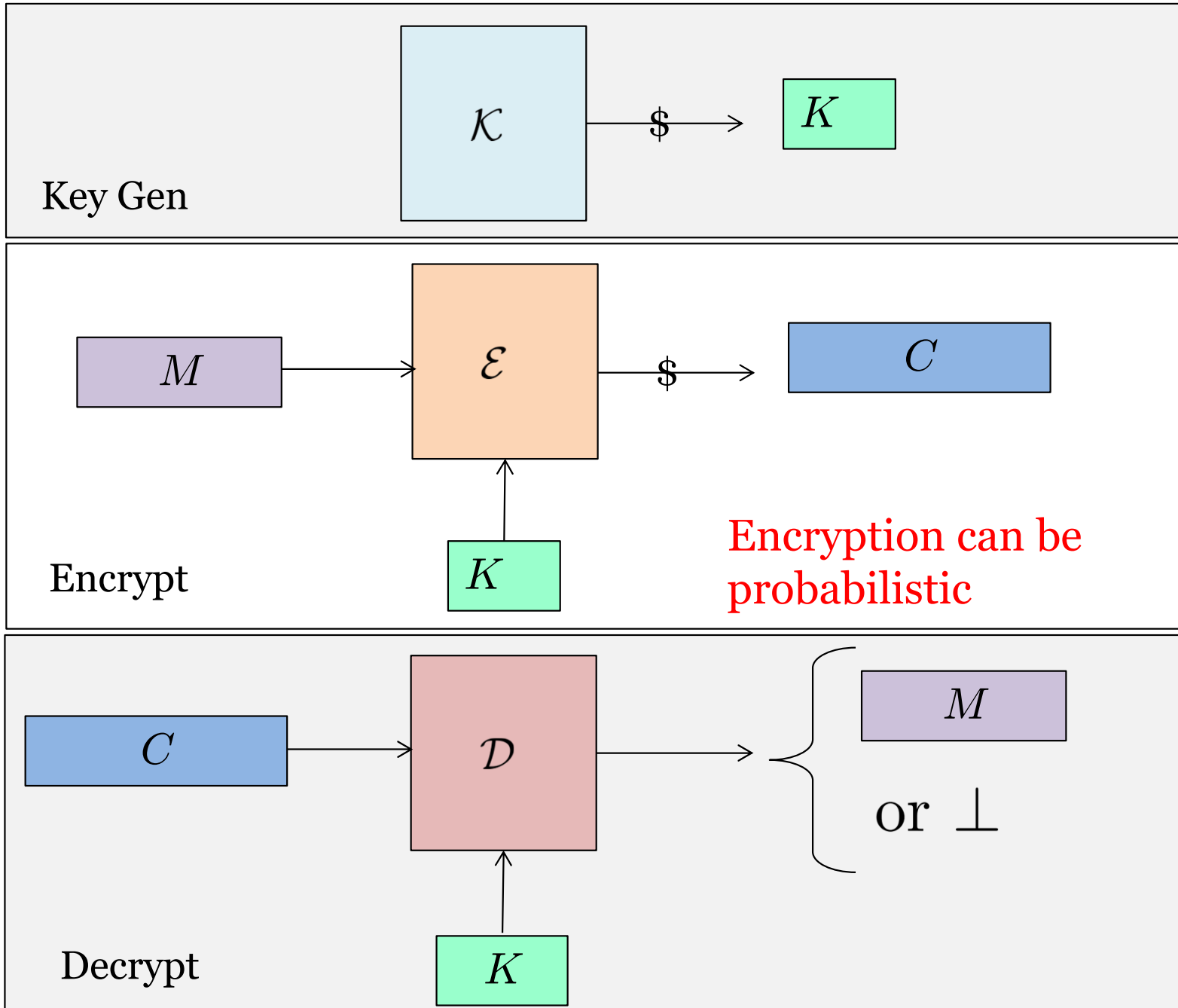
Agenda

1. Modes of Encryption: ECB, CBC, CTR

2. Formalizing Security

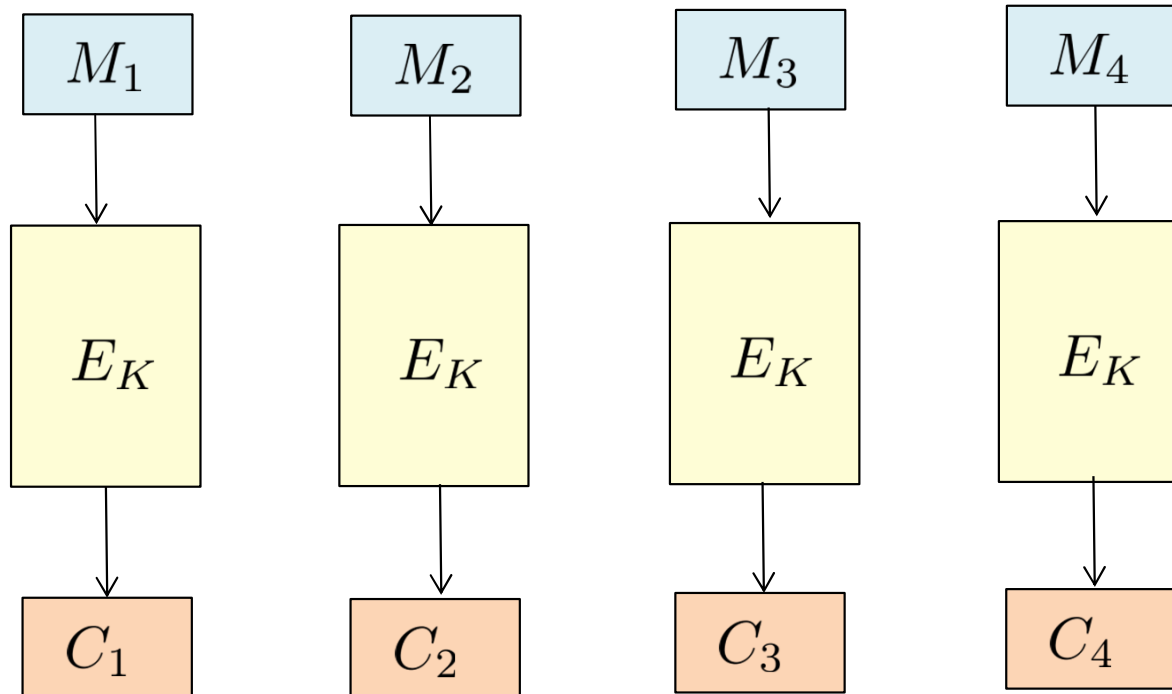
3. Stream Ciphers

Encryption Syntax



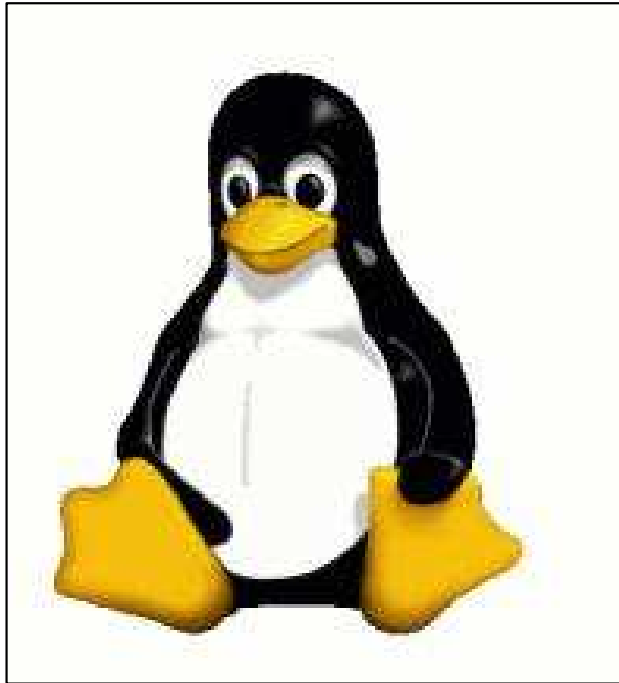
(Bad) Encryption Using Blockcipher: ECB

$$E : \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$$

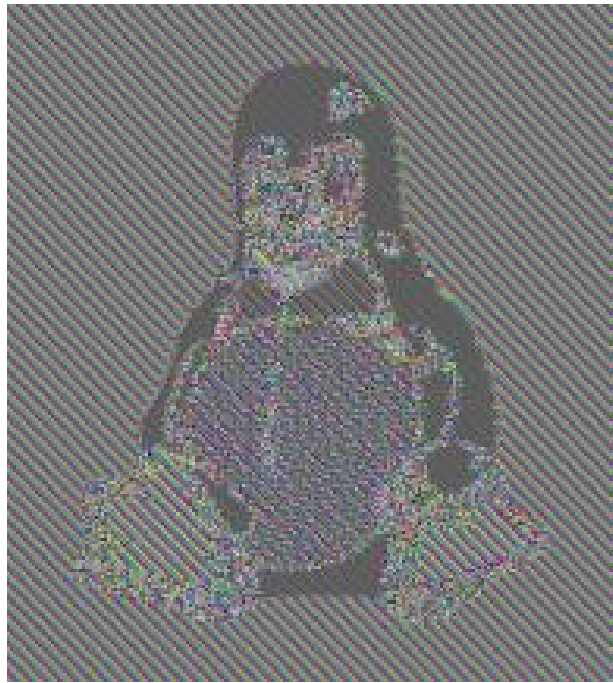


Can encrypt any message whose length is a multiple of n

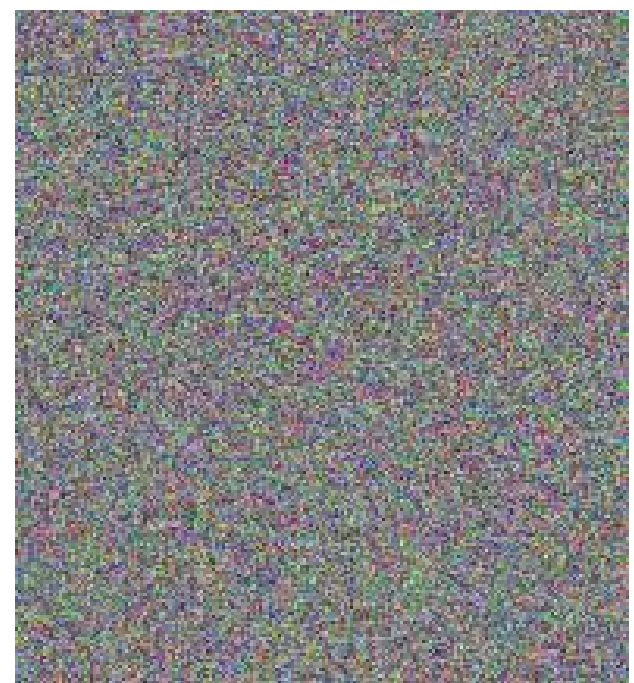
ECB Is **Insecure**



Message

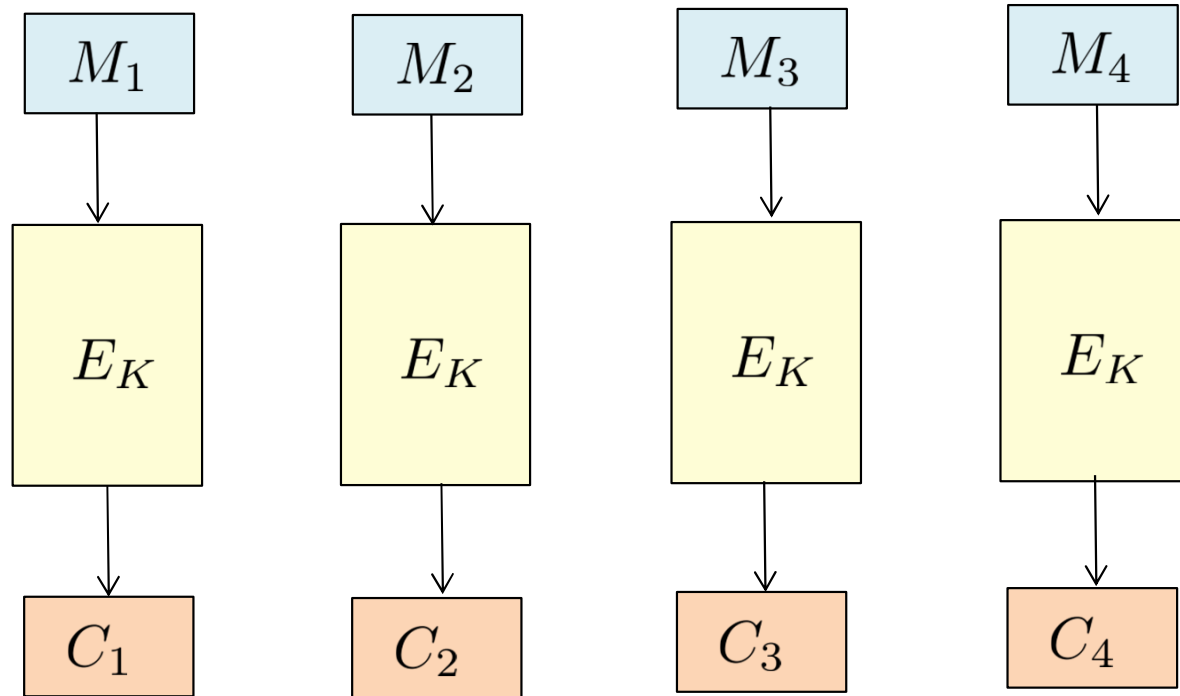


ECB ciphertext



Properly encrypted
ciphertext

Why Is ECB So Bad?



If $M_i = M_j$ then $C_i = C_j$

ECB Horror Stories

Half the apps in Android used ECB to encrypt data

An Empirical Study of Cryptographic Misuse in Android Applications

 ars TECHNICA

BIZ & IT —

How an epic blunder by Adobe
could strengthen hand of password
crackers

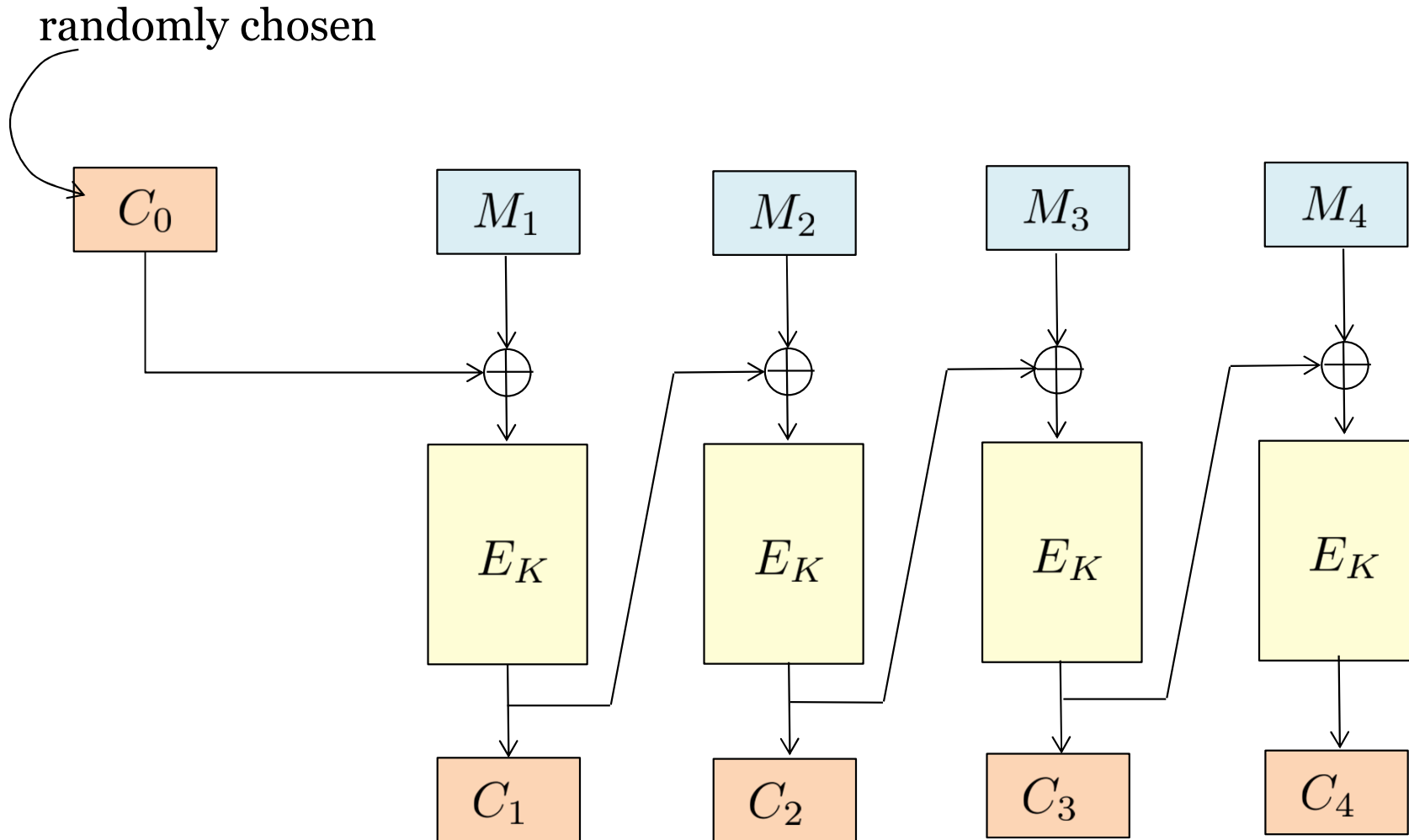
Adobe used ECB to
encrypt passwords

**Zoom concedes custom encryption is
substandard as Citizen Lab pokes holes in it**

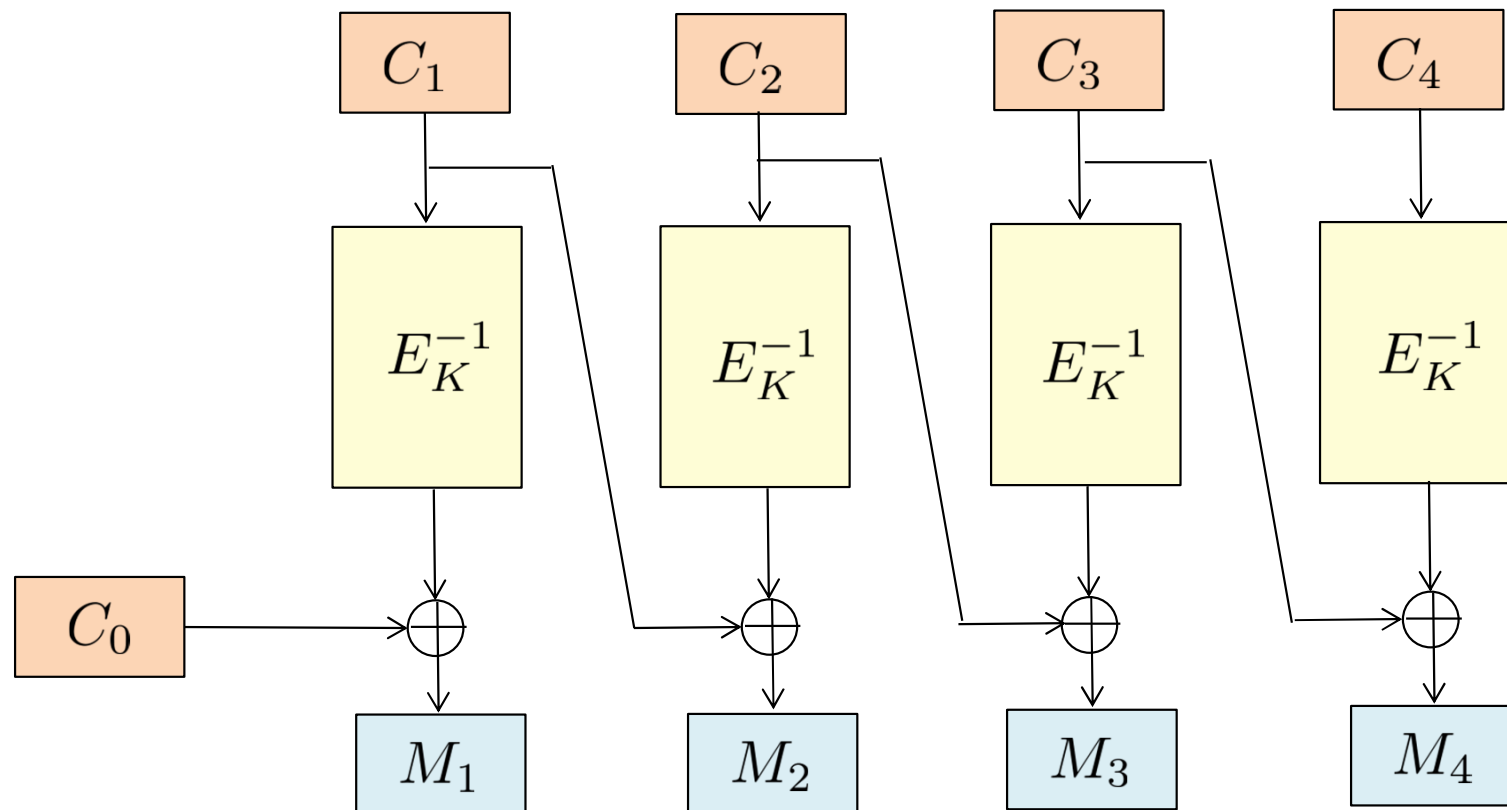
Zoom used ECB to encrypt video conferencing

Randomized Encryption: CBC

sequential



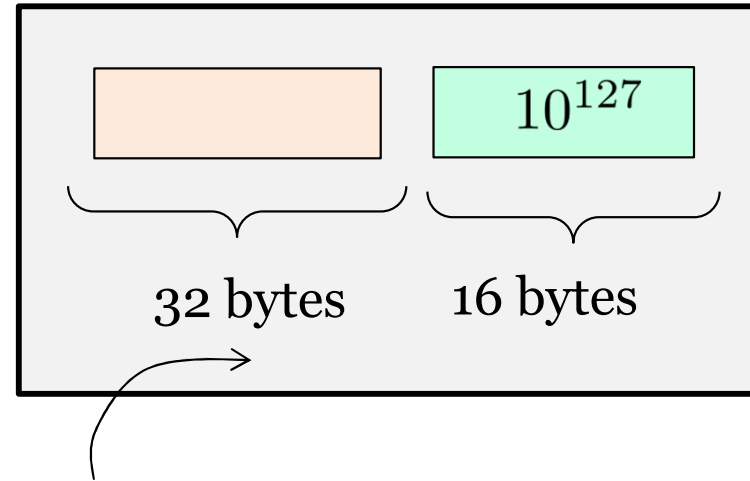
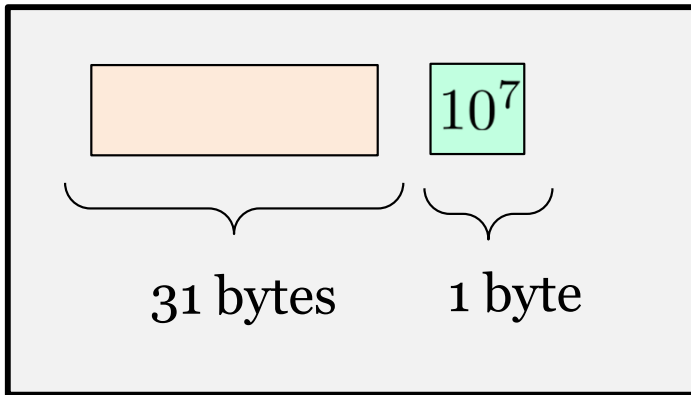
Decryption of CBC



Dealing with Fragmentary Data

Naive solution: Pad with 10^*

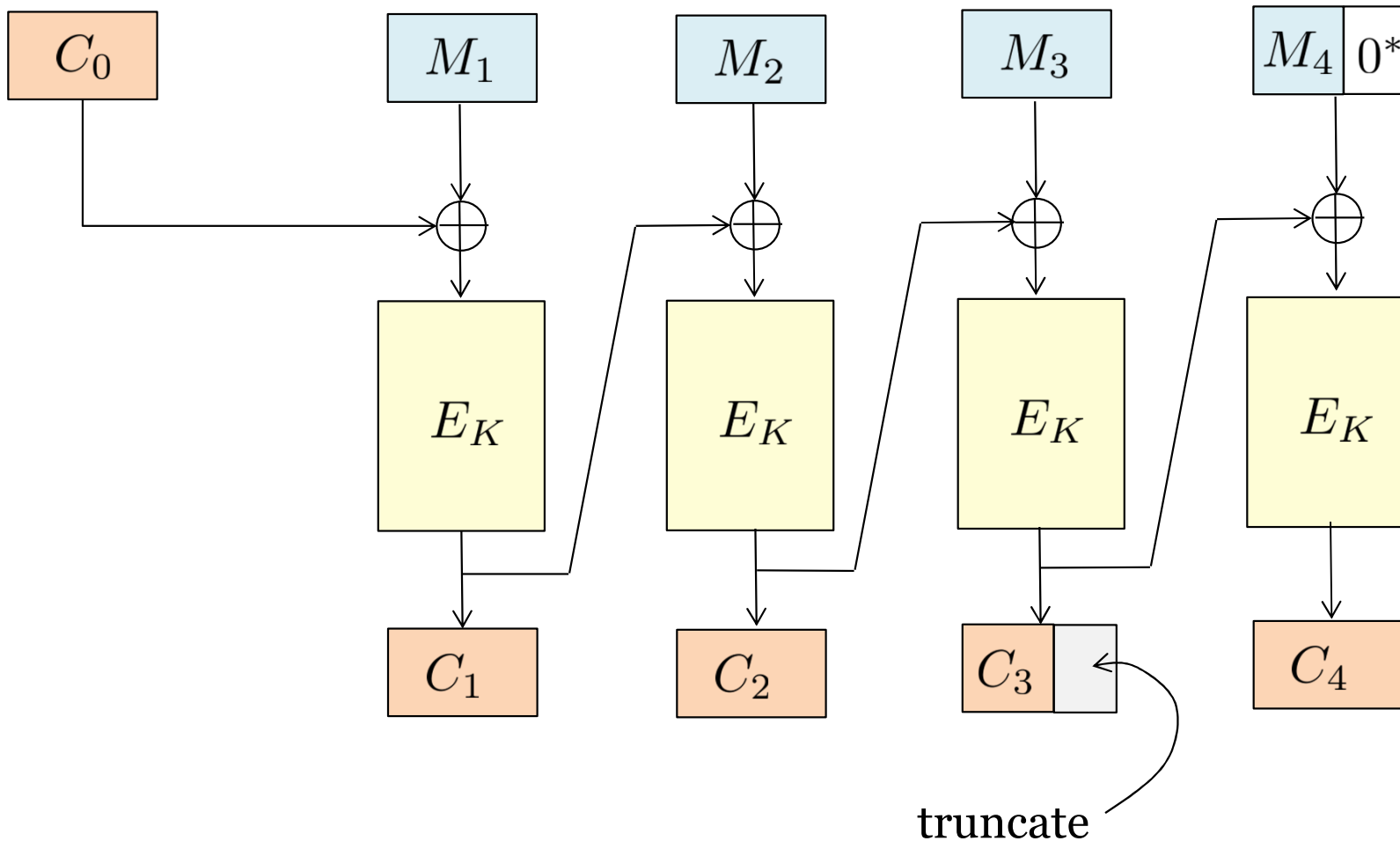
Example: Suppose that the block length is 16 bytes.



Padding is required, otherwise can't decrypt

Problem: Waste bandwidth, and for full-length msg, waste a blockcipher call

Ciphertext Stealing in CBC



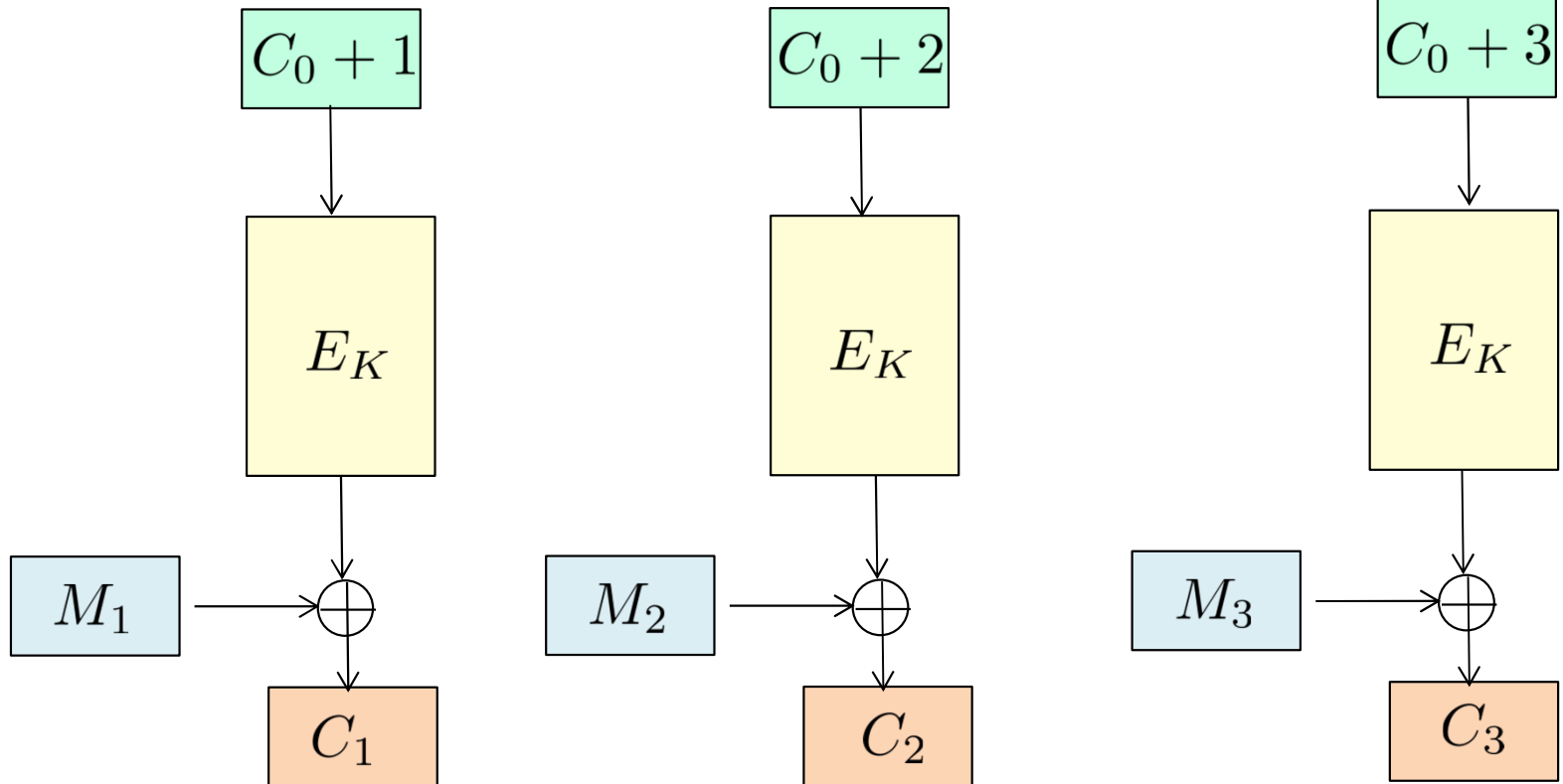
Exercise: How to use ciphertext stealing if msg is shorter than 1 block?

Randomized Encryption: CTR

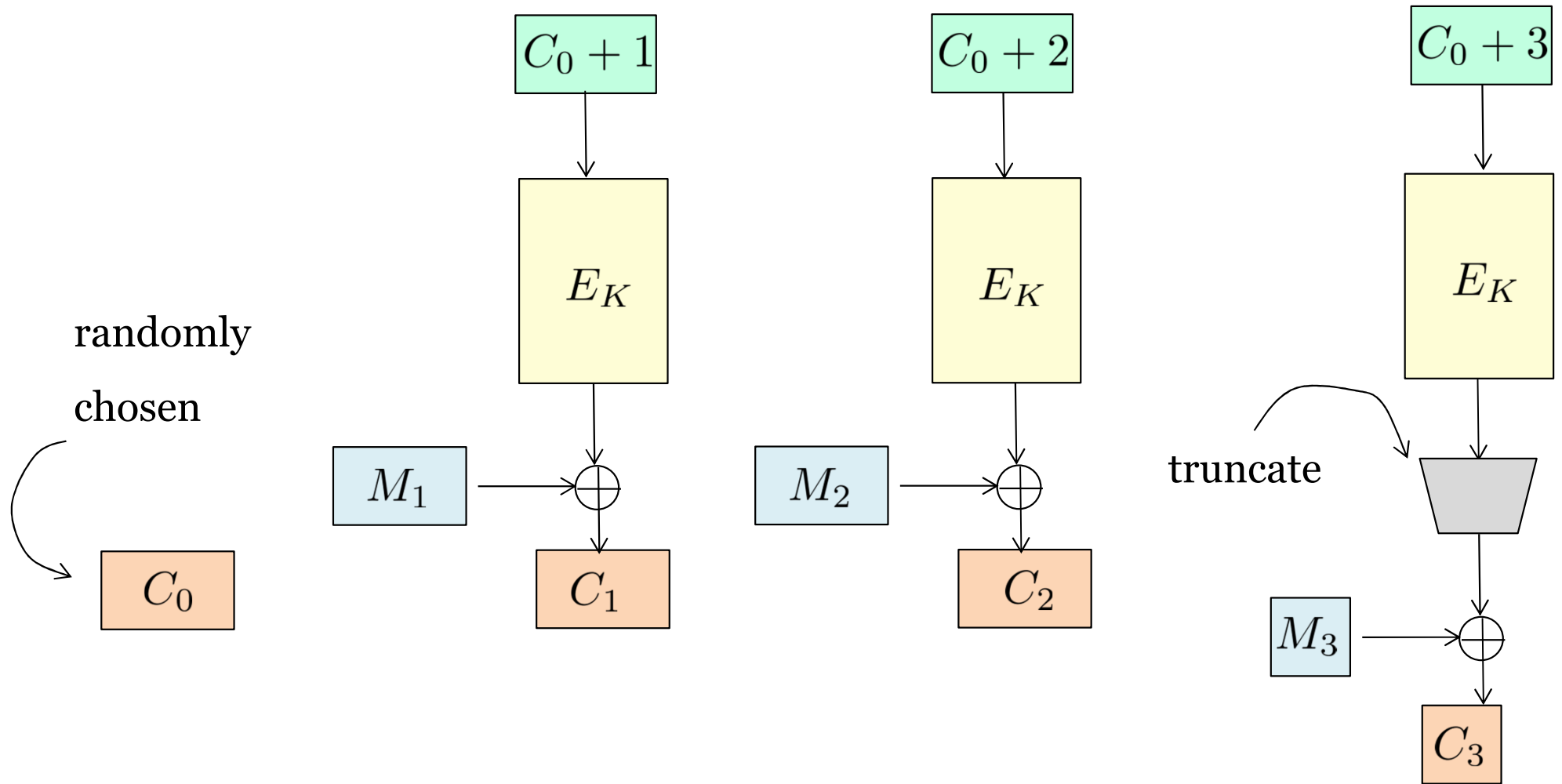
fully parallelizable

randomly
chosen

C_0



Dealing with Fragmentary Data



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3. Stream Ciphers



1982

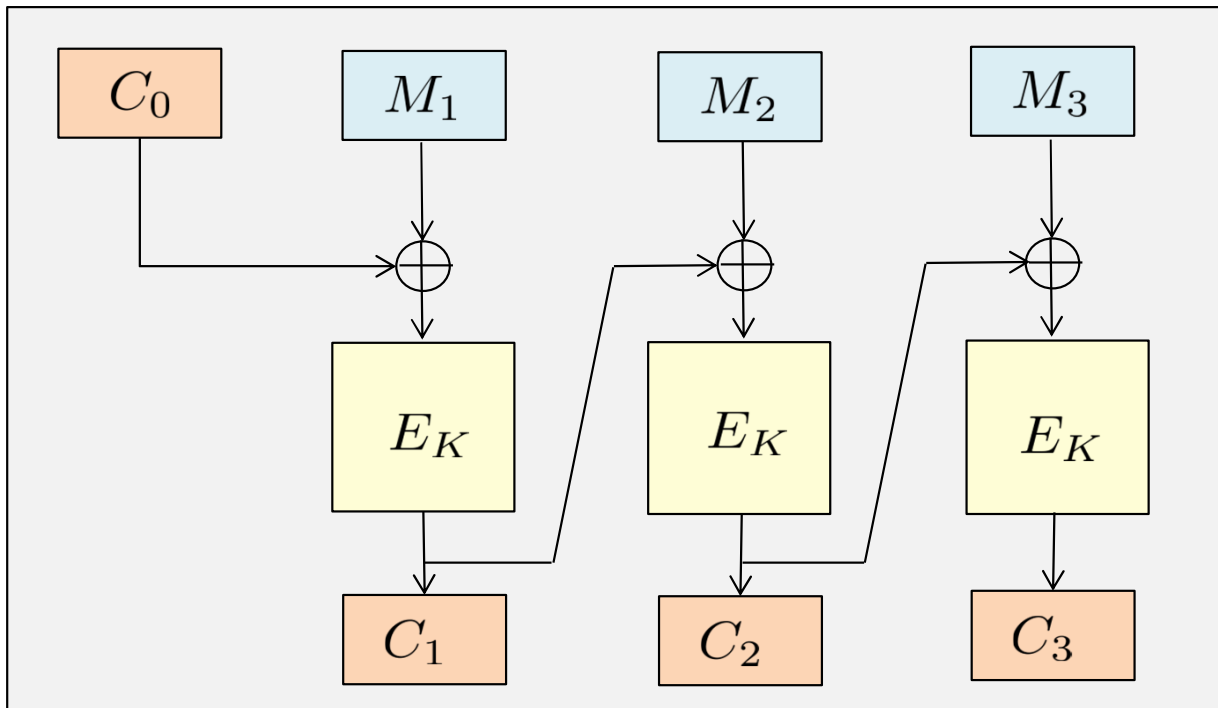
Formalizing Security: Intuition

Should hide

all partial information

about the plaintexts

Except message length



CBC trivially leaks
message length

Formalizing Security: Informal Definition

Adversary can't even distinguish the encryption of its **own chosen messages**

“A good disguise should not allow a mother to distinguish her own children”

Goldwasser and Micali

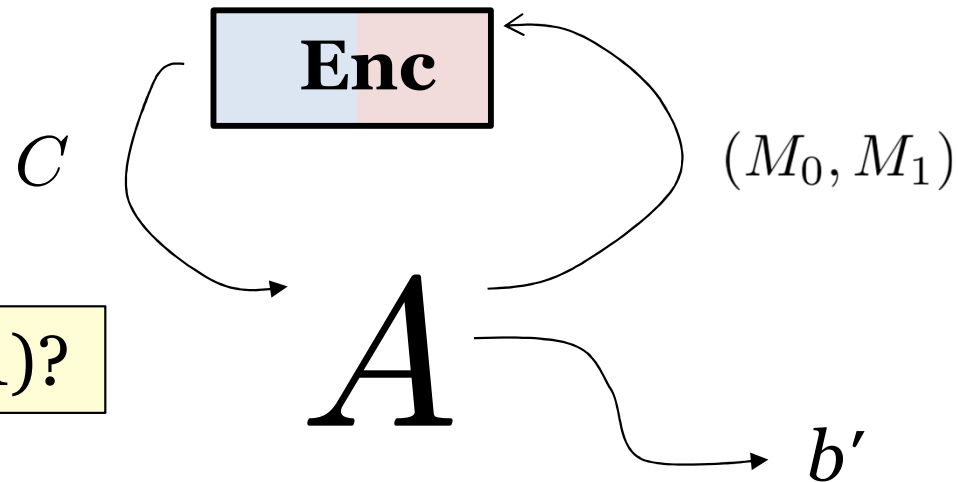
Formalizing Security: Left-or-Right

Left _{\mathcal{E}}

procedure **Enc**(M_0, M_1)
Return $\mathcal{E}_K(M_0)$

Right _{\mathcal{E}}

procedure **Enc**(M_0, M_1)
Return $\mathcal{E}_K(M_1)$



Left (0) or Right (1)?

$$\text{Adv}_{\mathcal{E}}^{\text{lr}}(A) = \Pr[\text{Right}_{\mathcal{E}}^A \Rightarrow 1] - \Pr[\text{Left}_{\mathcal{E}}^A \Rightarrow 1]$$

In each query, the two messages must have the same length

Formalizing Security: Real-or-Random

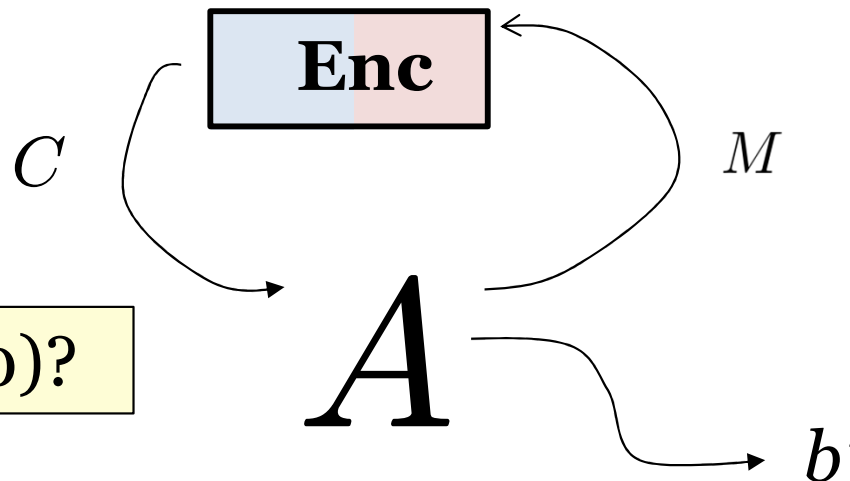
Real _{\mathcal{E}}

procedure **Enc**(M)
Return $\mathcal{E}_K(M)$

Rand _{\mathcal{E}}

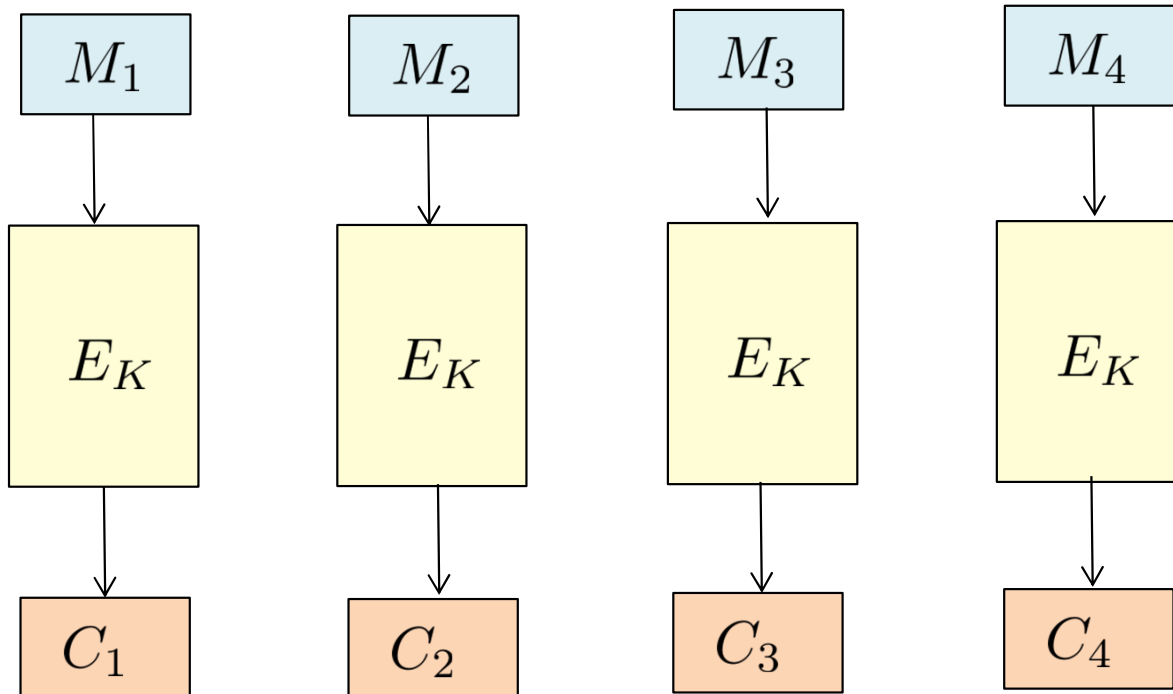
procedure **Enc**(M)
 $C \leftarrow \$ \mathcal{E}_K(M')$; $C' \leftarrow \$ \{0, 1\}^{|C|}$; Return C'

Real (1) or Rand (0)?

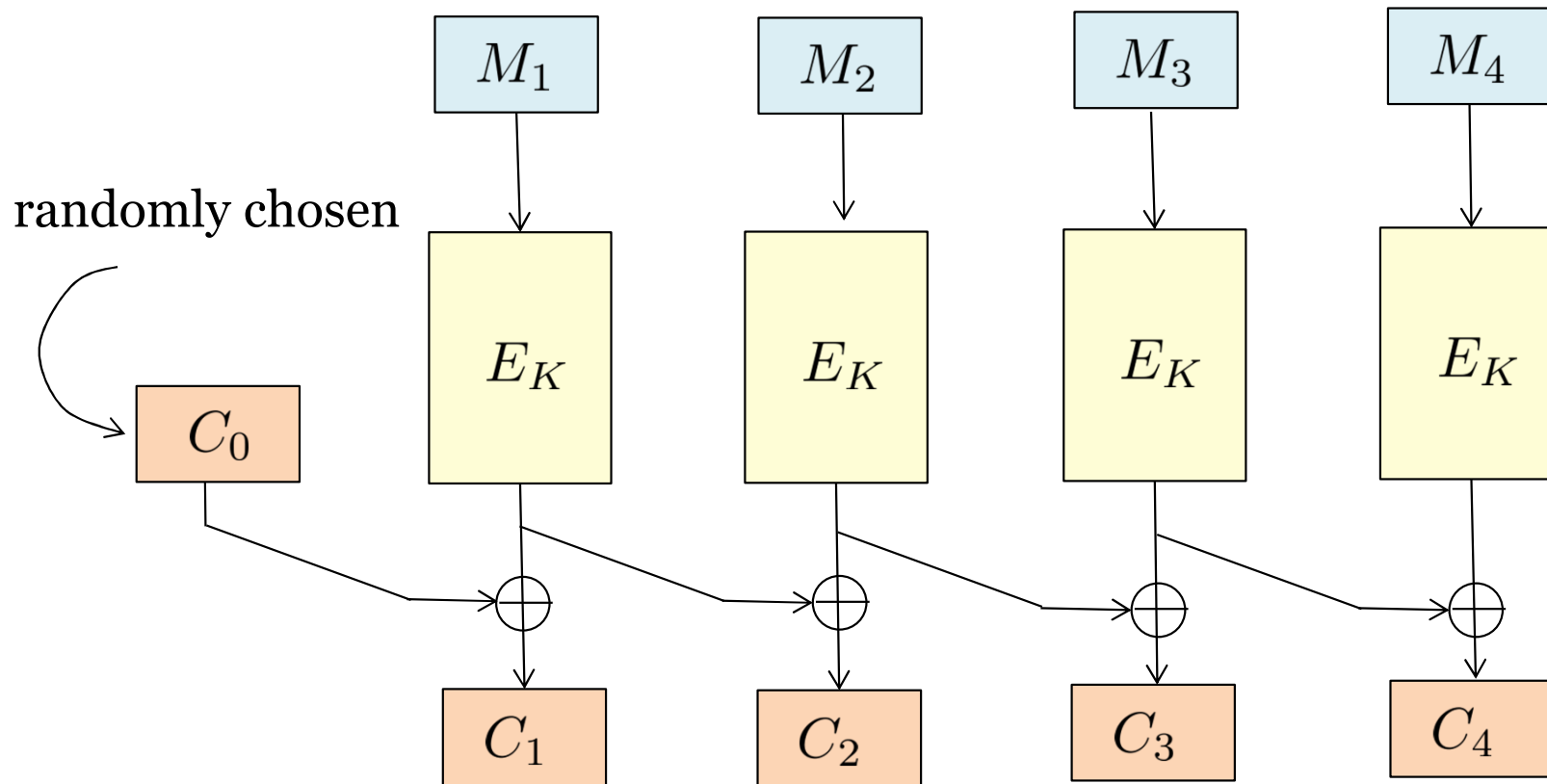


$$\text{Adv}_{\mathcal{E}}^{\text{rr}}(A) = \Pr[\text{Real}_{\mathcal{E}}^A \Rightarrow 1] - \Pr[\text{Rand}_{\mathcal{E}}^A \Rightarrow 1]$$

Exercise: Break LR Security of ECB



Exercise: Breaking RR Security



Question: Break the real-or-random security of this scheme using a single query of a 2-block message.

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Real-world Broken Stream Ciphers

RC4

Encryption scheme for Web traffic and Wifi traffic

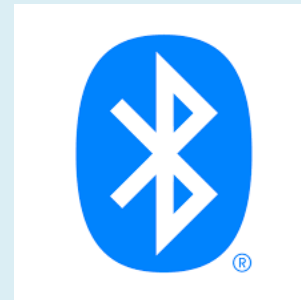


A5/1 Stream cipher



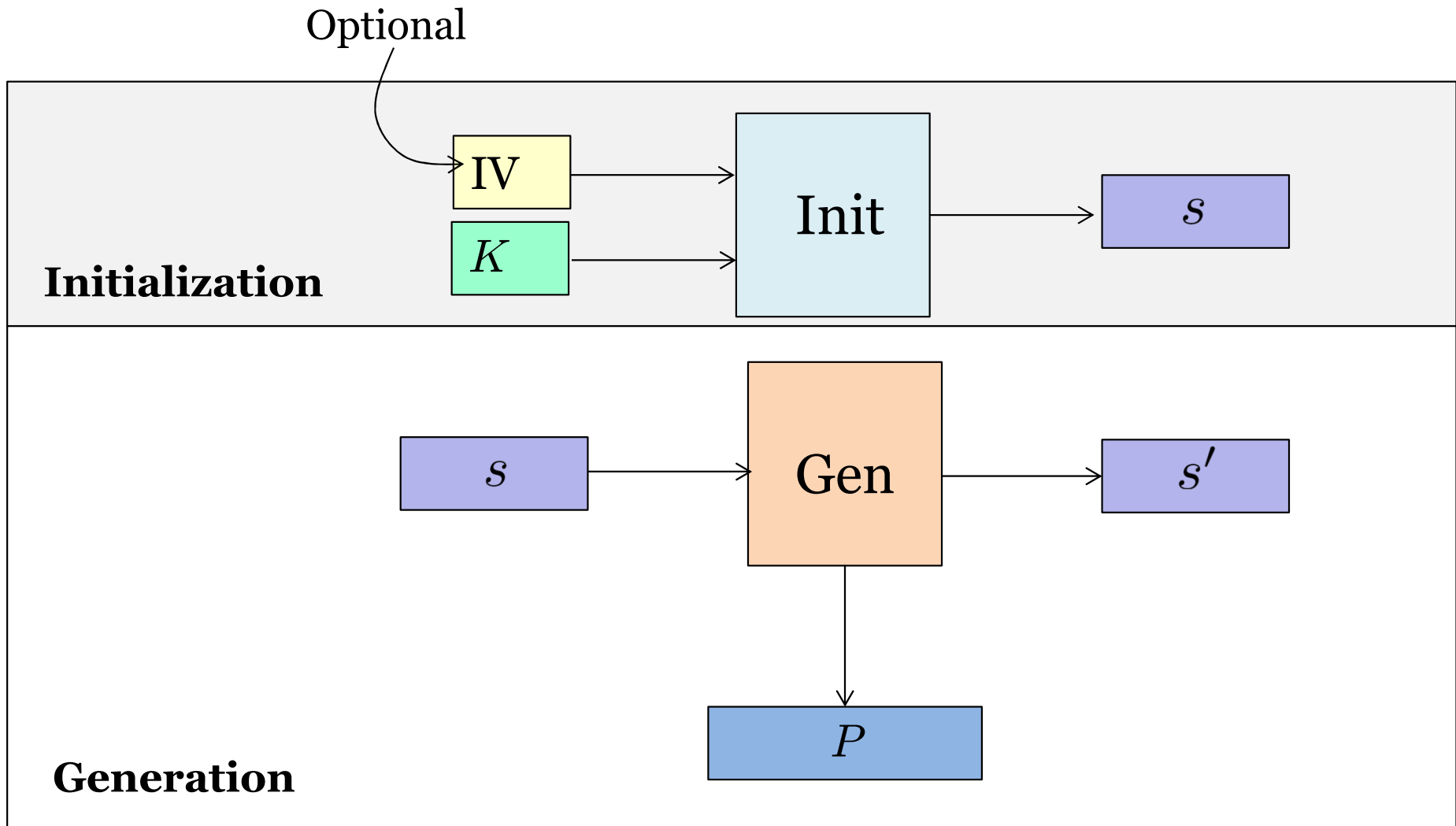
Cellular encryption
in GSM networks

E0 Stream cipher



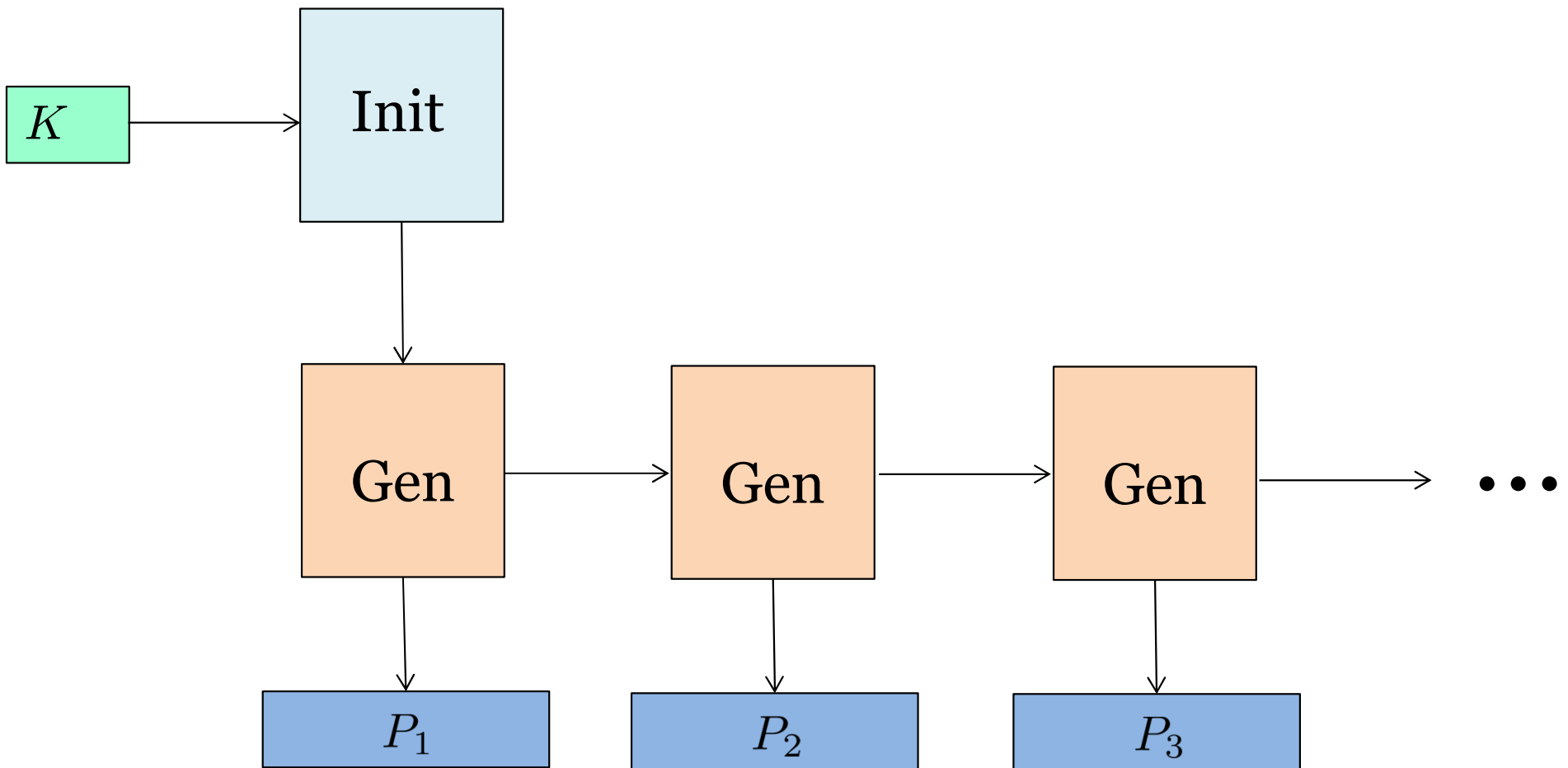
Bluetooth
encryption

Syntax



Use of Stream Cipher

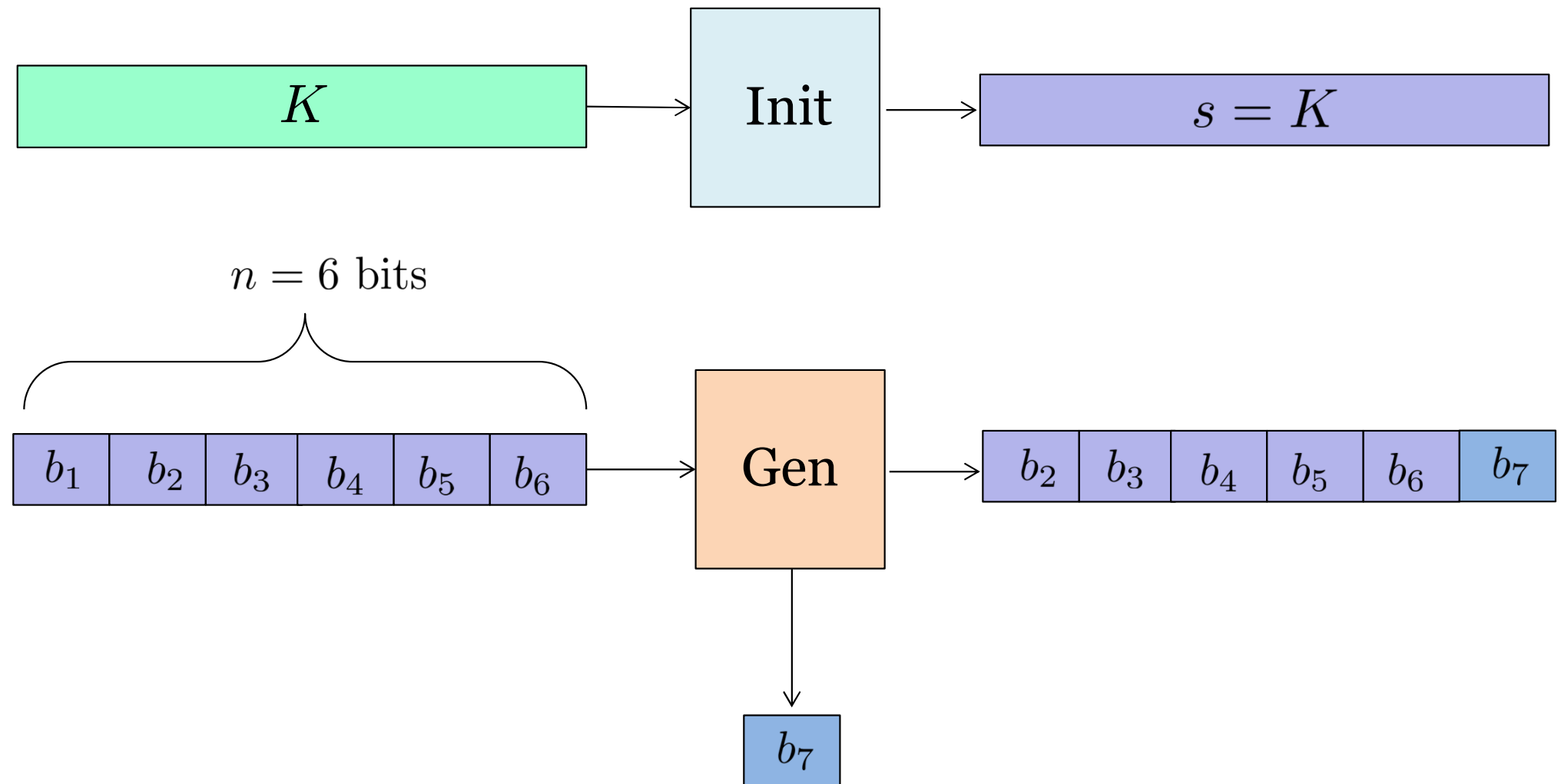
Producing A Stream of One-Time Pad



Question: Formalize a security notion for stream cipher

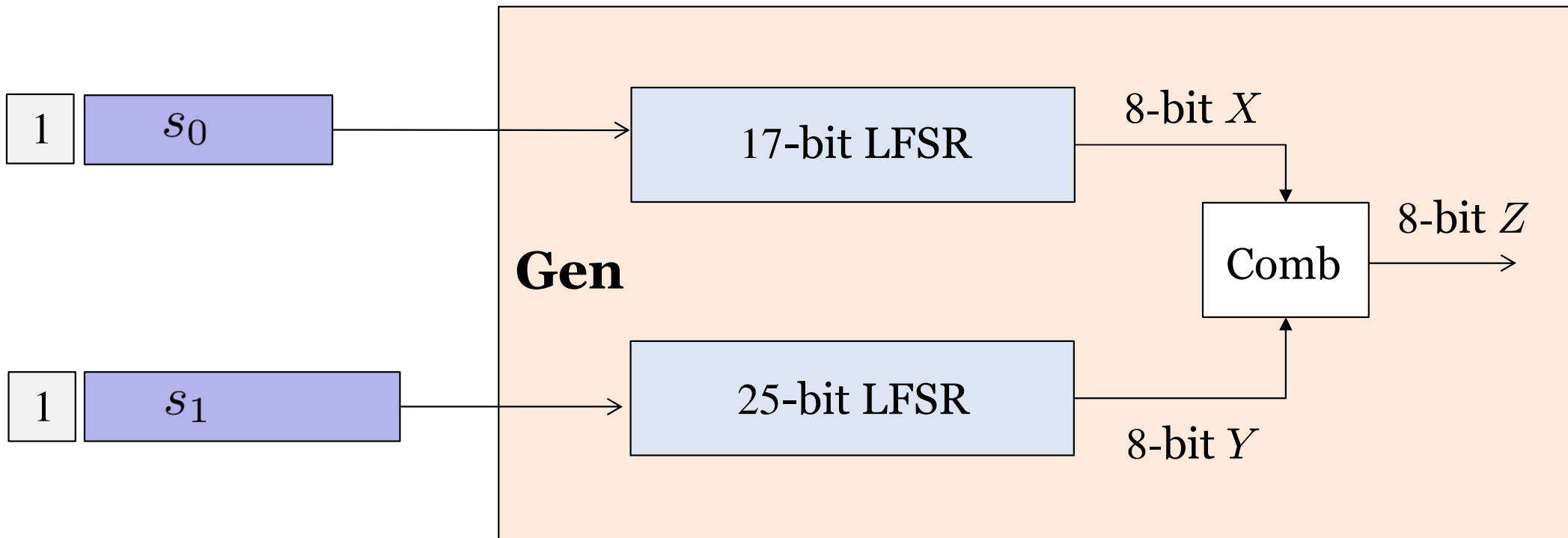
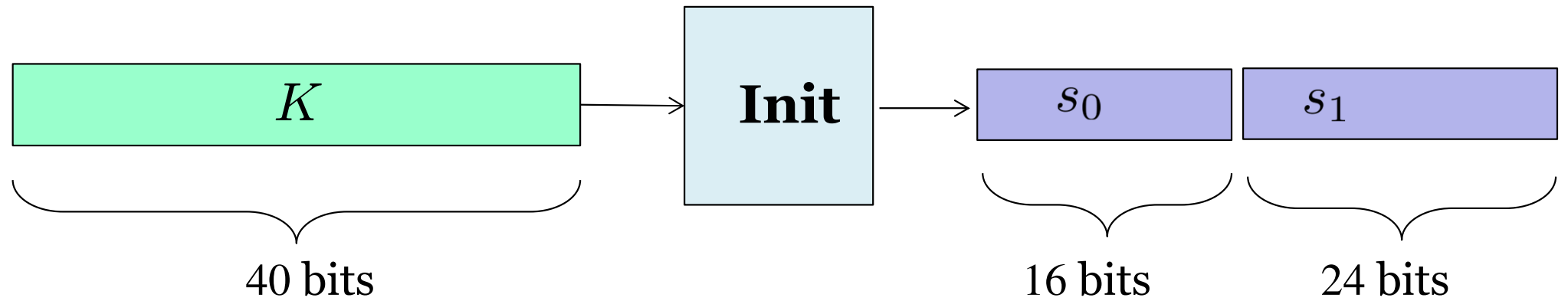
A Wrong Construction

Linear Feedback Shift Register (LFSR)

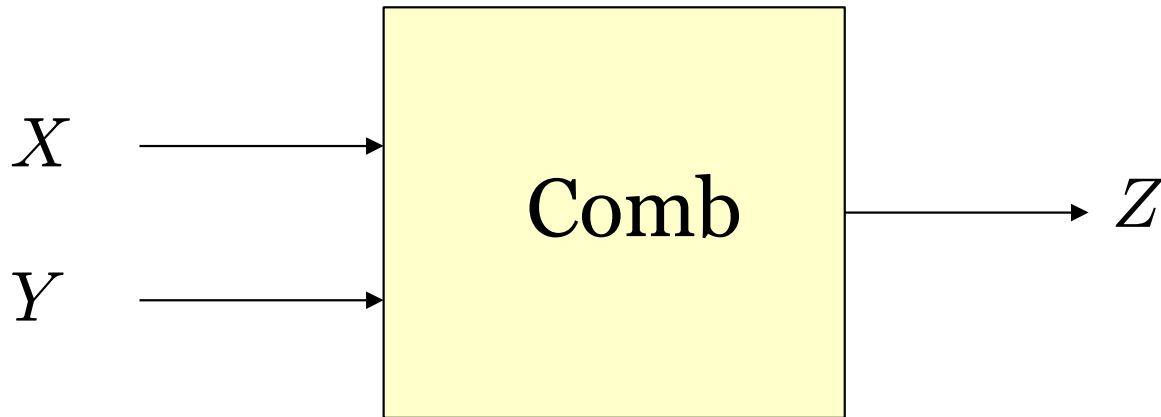


Question: Given n bits of output, recover subsequent bits

Case Study: DVD Encryption System

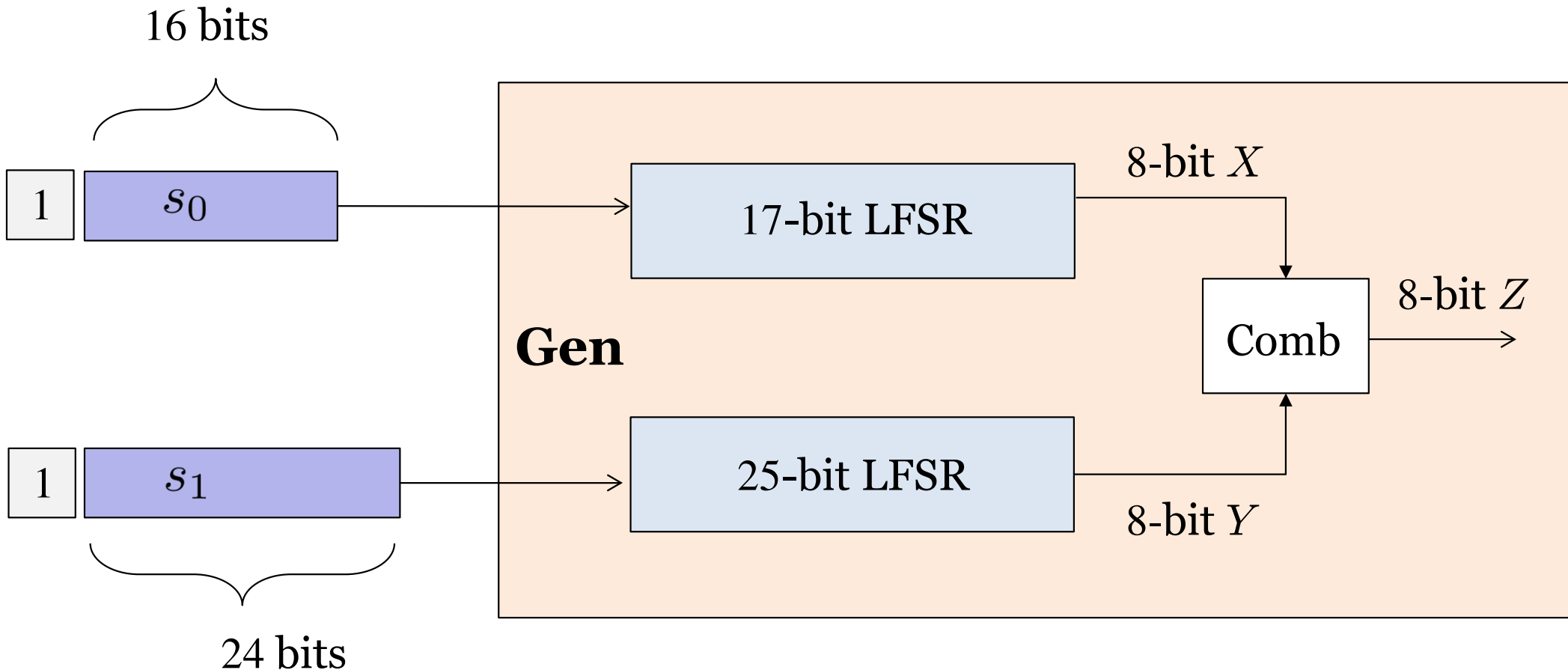


Property of Combiner To Exploit



Invertibility: Given Z and X , it's trivial to compute Y

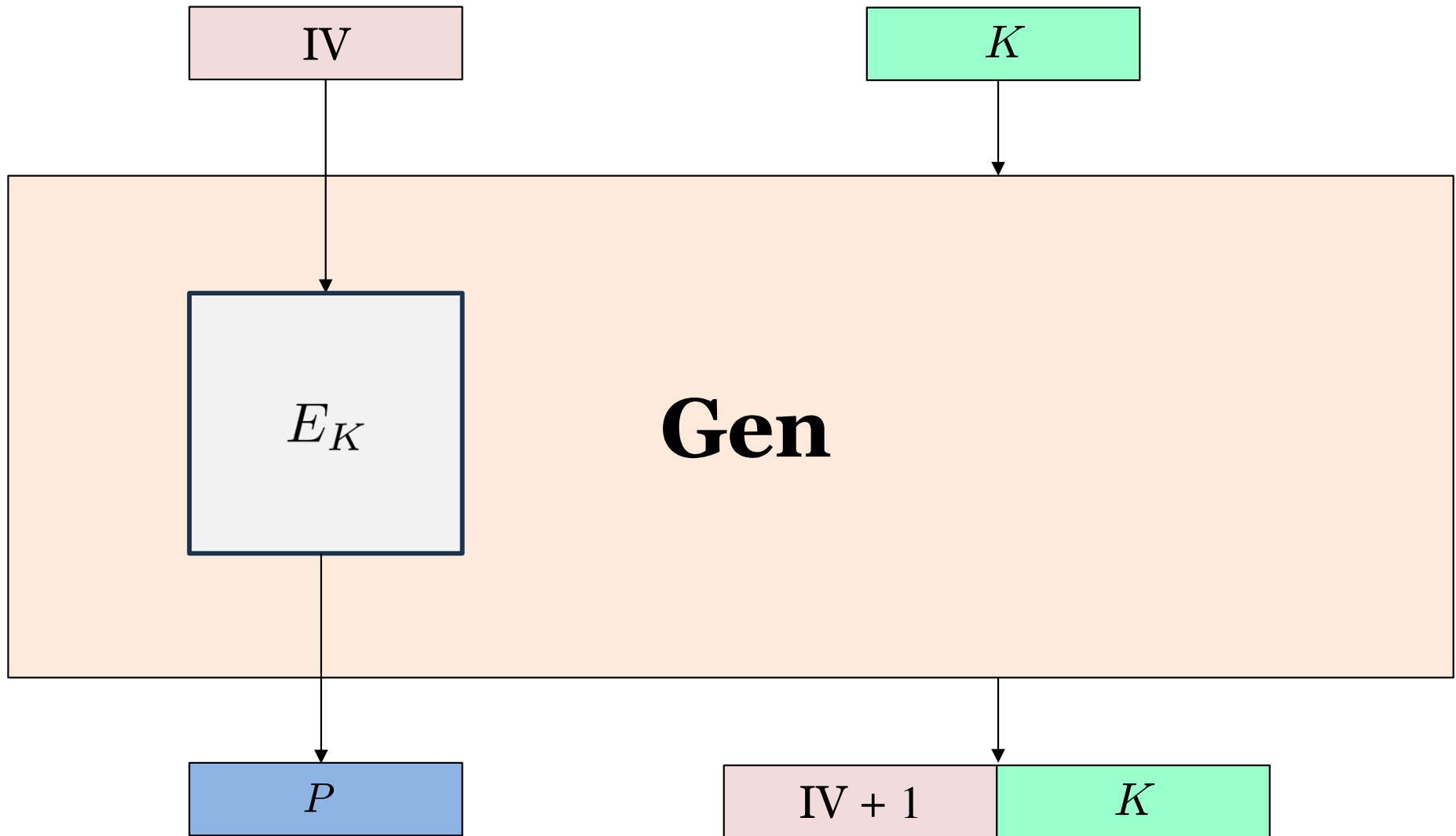
Breaking DVD Encryption System



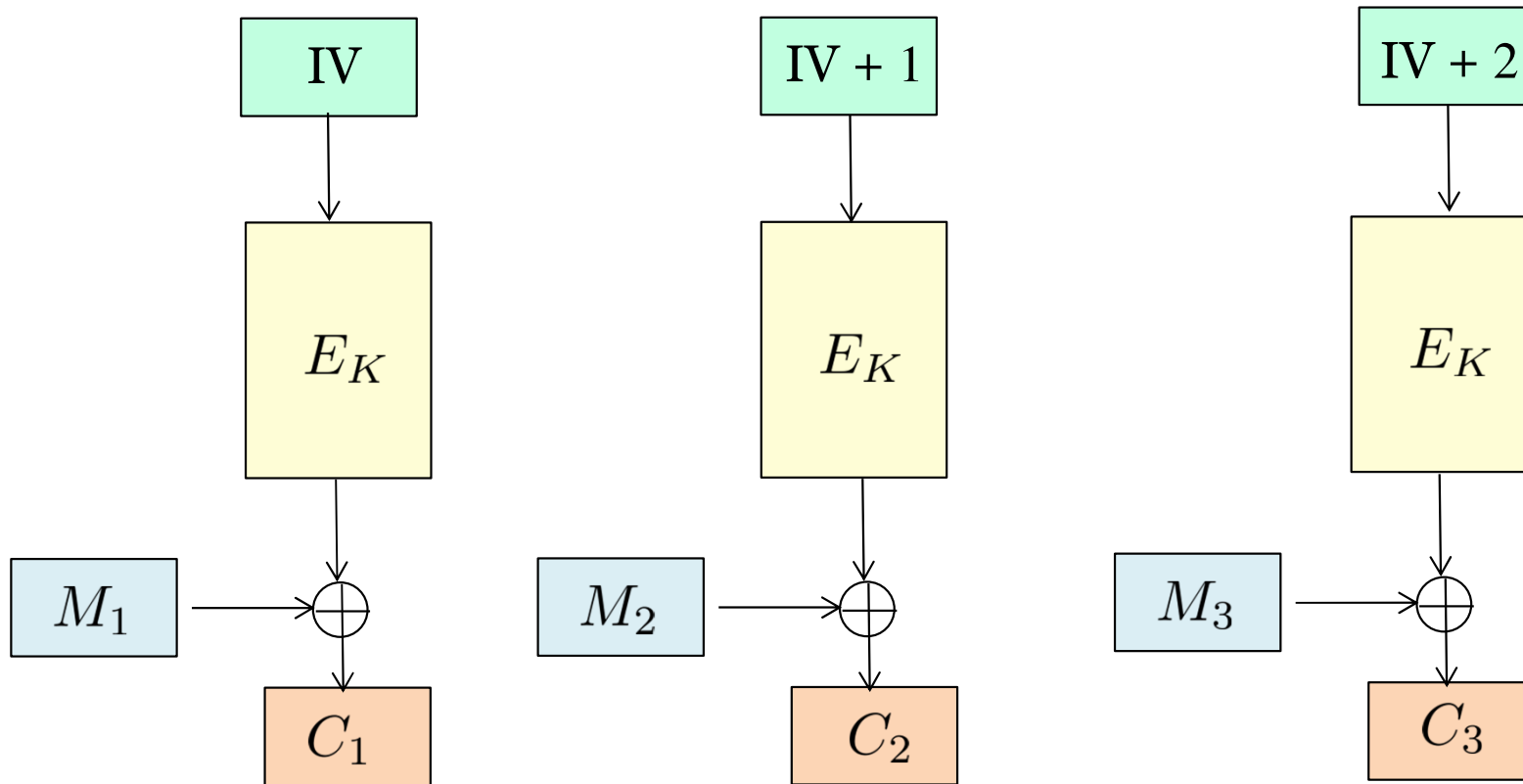
Question: Given the first 128 bits of output, recover subsequent bits using $O(2^{16})$ time by guessing the initial s_0

Building Stream Cipher From Blockcipher

Init sets $IV = 0$ and outputs (IV, K) as the initial state



How Encryption Looks Like: Stateful CTR



Ciphertext doesn't include IV

Sender and receiver update $IV \leftarrow IV + 3$ for the next encryption